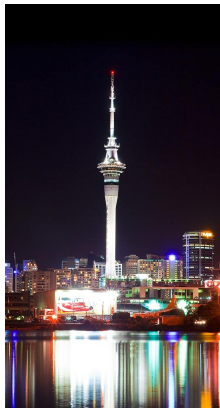


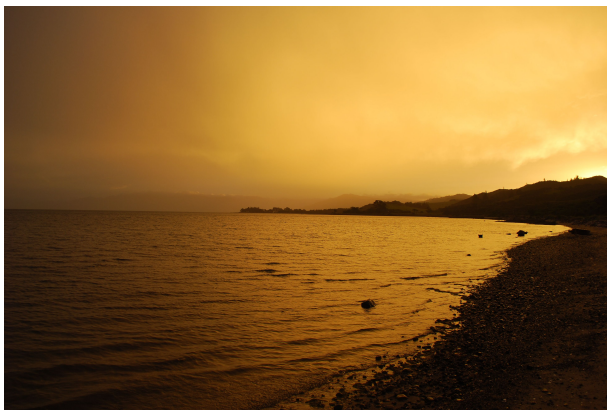
# Distinguishing Maximal Orders of Quaternion Algebras by their Short Elements

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Speakers: Pierre Deligne, Gus Lehrer, Cheryl Praeger, René Schoof,  
Richard Weiss.



Auckland, New Zealand, December 2015

# Plan

- ▶ Background and some of my favourite questions
- ▶ Why is this talk in a session on lattices?
- ▶ Sketch of results and algorithm
- ▶ Shall we talk about something else?

Thanks: David Kohel, Drew Sutherland.

Please ask questions at any time.

Ilya Chevyrev



# Elliptic Curves and Isogenies

- ▶ An **elliptic curve** over a field  $\mathbb{k}$  is a non-singular projective cubic curve. The set of  $\mathbb{k}$ -rational points is a group.
- ▶ An **isogeny**  $\phi : E_1 \rightarrow E_2$  of elliptic curves is a morphism that is a group homomorphism.
- ▶ Isogenies satisfy a degree 2 characteristic polynomial  $T^2 - \text{Tr}(\phi)T + \text{deg}(\phi) = 0$ , having discriminant  $D = \text{Tr}(\phi)^2 - 4 \text{deg}(\phi) \leq 0$ .
- ▶ Tate's isogeny theorem: Let  $E_1, E_2$  be elliptic curves over a finite field  $\mathbb{F}_q$ . Then  $\#E_1(\mathbb{F}_q) = \#E_2(\mathbb{F}_q)$  iff there is an isogeny  $\phi : E_1 \rightarrow E_2$  over  $\mathbb{F}_q$ .
- ▶  $\text{End}(E) = \{\text{isogenies } \phi : E \rightarrow E \text{ over } \overline{\mathbb{F}_q}\}$ .
- ▶  $\text{End}(E)$  is either an order in an imaginary quadratic field (ordinary) or a maximal order in a definite quaternion algebra  $B_p$  ramified at  $\{p, \infty\}$  (supersingular).

## Some Computational Questions

- ▶ Given  $E$  over  $\mathbb{F}_q$  to compute  $\text{End}(E)$ .  
Two cases: ordinary and supersingular.
- ▶ Given  $E, E'$  over  $\mathbb{F}_q$  with  $\#E(\mathbb{F}_q) = \#E'(\mathbb{F}_q)$  to compute an isogeny from  $E'$  to  $E$ .  
Two cases: ordinary and supersingular.
- ▶ Given  $q, N$  construct an elliptic curve  $E/\mathbb{F}_q$  with  $\#E(\mathbb{F}_q) = N$ .
- ▶ Construct an elliptic curve  $E/\mathbb{F}_q$  with  $\#E(\mathbb{F}_q) = N$  for pairs  $(q, N)$  with certain “desired properties”.
- ▶ Given a maximal order  $\mathcal{O}$  in the quaternion algebra  $B_p$  to construct an elliptic curve  $E$  over  $\mathbb{F}_p$  or  $\mathbb{F}_{p^2}$  with  $\text{End}(E) \cong \mathcal{O}$ .

# Hilbert Class Polynomial

- ▶ Consider fundamental discriminant  $D < 0$ . The Hilbert class polynomial  $H_D(X) \in \mathbb{Z}[X]$  has property:  
Given  $E/\mathbb{k}$ , if  $H_D(j(E)) = 0$  then  $\text{End}(E)$  contains an isogeny of discriminant  $D$ .
- ▶ Specifically, the roots in  $\mathbb{C}$  of  $H_D(X)$  are the  $j$ -invariants of the elliptic curves over  $\mathbb{C}$  possessing the quadratic order  $\mathcal{O}_D = \mathbb{Z}[\frac{1}{2}(D + \sqrt{D})]$  as their endomorphism ring.
- ▶ Class polynomials are used in the CM method for constructing curves with a given group order/endomorphism structure.
- ▶ What other applications might there be?



# Bröker's Algorithm

- ▶ Goal: Given  $q = p^a$  to construct a supersingular curve over  $\mathbb{F}_q$  with specified trace of Frobenius.
- ▶ Main idea: Choose small prime  $\ell$  such that  $\left(\frac{-\ell}{p}\right) = -1$  then find root of  $H_{-\ell}(X)$  or  $H_{-4\ell}(X)$  in  $\mathbb{F}_q$ .
- ▶ Construct corresponding  $E$  and twist if necessary.
- ▶ CM theory tells that  $E$  is supersingular, as  $p$  is inert in  $\mathbb{Q}(\sqrt{-\ell})$ .

## Idea

- ▶ Problem: Given a maximal order  $\mathcal{O}$  construct  $E$  such that  $\text{End}(E) \cong \mathcal{O}$ .
- ▶ A simple idea is to find some elements in  $\mathcal{O}$  of small discriminant  $D_1, D_2, \dots$  and take

$$G(X) = \gcd(H_{D_1}(X), H_{D_2}(X), \dots).$$

- ▶ Then hope that  $\deg(G) \leq 2$  and that taking roots gives  $j(E)$  and hence  $E$ .
- ▶ Related application: Can we determine  $\text{End}(E)$  by testing if  $H_D(j(E)) = 0$  for various discriminants  $D$ ?
- ▶ **Question:** Is a maximal order  $\mathcal{O}$  in quaternion algebra  $B_p$  determined by a small number of discriminants.

# Lattices and Ternary Forms

- ▶ Consider the  $\mathbb{Z}$ -module  $\mathcal{O}^T = \{2x - \text{Tr}(x) : x \in \mathcal{O}\}$  of rank 3.
- ▶ Note that  $y \in \mathcal{O}^T$  implies  $\text{Tr}(y) = 0$  (pure quaternion).
- ▶ The reduced norm on  $\mathcal{O}$  is a ternary quadratic form  $Q$ , making  $\mathcal{O}^T$  a lattice.
- ▶ The volume of the lattice is  $4p^2$ .
- ▶ Let  $\mathcal{O}'$  be another maximal order in the same quaternion algebra  $B_p$  and let  $Q'$  be the ternary form of  $\mathcal{O}'^T$ .  
If  $Q'$  is equivalent to  $Q$ , in the sense of quadratic forms, then is  $\mathcal{O}'$  isomorphic to  $\mathcal{O}$  ( $\mathcal{O}' = c\mathcal{O}c^{-1}$  for some  $c \in B_p$ )?
- ▶ **Theorem:** (Schiemann) Ternary quadratic forms are determined up to equivalence by their theta series.
- ▶ We will show that one can check equivalence by only checking a very small number of coefficients of the theta series.

# Bulguksa Lattices



# Main Theorems

## Theorem

Let  $\mathcal{O}$  and  $\mathcal{O}'$  be two maximal orders of  $B_p$ . Let  $\mathcal{O}^T$  and  $\mathcal{O}'^T$  have the same successive minima  $D_1 \leq D_2 \leq D_3$ . Assume moreover that  $D_1 D_2 < 16p/3$  and that  $p$  is sufficiently large. Then  $\mathcal{O}$  and  $\mathcal{O}'$  are of the same type (= isomorphic).

## Theorem

Let  $p > 286$  and  $\mathcal{O}, \mathcal{O}'$  be two maximal orders of  $B_p$ . Let  $D_1, D_2$  and  $D_3$  be the successive minima of  $\mathcal{O}^T$  and let  $x, y \in \mathcal{O}^T$  be such that  $Nr(x) = D_1$  and  $Nr(y) = D_2$ . Suppose that  $D_1 D_2 < \frac{16}{3}p$  and that  $D_1, D_2, Nr(x+y), Nr(x-y)$  and  $D_3$  are all "represented optimally" in  $\mathcal{O}'^T$  and that  $\theta'_{\mathcal{O}^T}(D_3) \leq \theta'_{\mathcal{O}'^T}(D_3)$ . Then  $\mathcal{O}$  and  $\mathcal{O}'$  are of the same type.

The condition  $D_1 D_2 < 16p/3$

### Lemma

*Let  $\mathcal{O}$  be a maximal order in  $B_p$  that contains an element  $\pi$  such that  $\pi^2 = -p$  (and hence  $j(\mathcal{O}) \in \mathbb{F}_p$ ). Then  $D_1 D_2 < 16p/3$ .*

Proof based on a paper of Kaneko.

Elkies showed  $D_1 \leq 2p^{2/3}$  for any maximal order in  $B_p$  and Yang has shown that this is best possible.

## Method of Proof

- ▶ Let  $x, y \in \mathcal{O}^T$  have norms  $D_1$  and  $D_2$  respectively. Similarly  $x', y' \in \mathcal{O}'^T$ .
- ▶ Prove that  $\langle x, y \rangle$  and  $\langle x', y' \rangle$  isometric, using  $4D_1D_2 - \text{Tr}(x\bar{y})^2 \equiv 0 \pmod{p}$  and simple geometry of numbers.
- ▶ **Lemma:**  $w = 2xy - \text{Tr}(xy) \in \mathcal{O}^T \cap \langle x, y \rangle^\perp$ .
- ▶ More geometry of numbers completes the result.
- ▶ Proof of Theorem 2 requires further arguments to reduce to case of Theorem 1.
- ▶ Everyone agrees there should be a nicer proof.

## Algorithm to Construct $E$

- ▶ Let  $\mathcal{O}$  be a maximal order in  $B_p$  given as a  $\mathbb{Z}$ -basis.
- ▶ Use lattice algorithms to find several small norms  $d_1, d_2, \dots, d_n$  of “primitive” elements in  $\mathcal{O}^T$ .
- ▶ Hence  $(X - j(E))$  is a factor of  $\gcd(H_{-d_1}(X), H_{-d_2}(X), \dots, H_{-d_n}(X))$ .
- ▶ Take multiple roots into account.
- ▶ When  $j(E) \in \mathbb{F}_p$  then our theorems imply the algorithm terminates with a degree 1 polynomial.
- ▶ In this case, all  $d_i$  are such that  $|d_i| = O(p)$ .
- ▶ Computing  $H_d(X)$  can be done in time  $\tilde{O}(|d|)$  by Belding-Bröker-Enge-Lauter or Sutherland.  
This is the limiting step, as poly degree is  $O(|d|^{0.5+\epsilon})$ .
- ▶ So overall complexity  $\tilde{O}(p)$ .
- ▶ Examples in paper.



## Algorithm when $j(E) \notin \mathbb{F}_p$

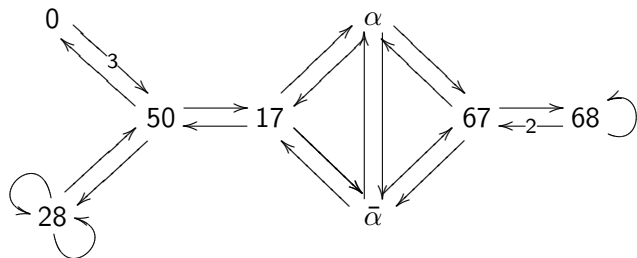
- ▶ Conjecture that the algorithm terminates with degree two polynomial.
- ▶ Conjecture that running time is still  $\tilde{O}(p)$ .
- ▶ Can consider an algorithm to match  $\{\mathcal{O}\}$  with the set  $\{j(E)\}$  over all supersingular curves.
- ▶ Cerviño proposed such an algorithm.  
As far as we can tell, his algorithm requires  $O(p^{3+\varepsilon})$  field operations.
- ▶ Our method has the improved complexity  $O(p^{2.5+\varepsilon})$  field operations.
- ▶ Our algorithm is always guaranteed to halt!
- ▶ For subcase of  $j(E) \in \mathbb{F}_p$ , Cerviño needs  $O(p^{2.5+\varepsilon})$  and we need  $O(p^{1.5+\varepsilon})$ .

The last talk of the conference has tools that should lead to better solutions to these problems.

# Computing Isogenies between Supersingular Elliptic Curves over $\mathbb{F}_p$

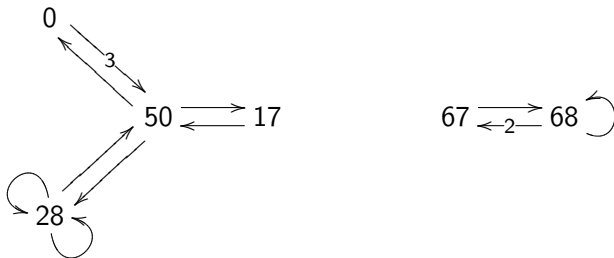
- ▶ Joint work with Christina Delfs.
- ▶ Problem is to find sequence of isogenies between two given supersingular elliptic curves.
- ▶ The number of supersingular elliptic curves in  $\overline{\mathbb{F}}_p$  is approximately  $p/12$ , but there are only  $p^{0.5+\epsilon}$  supersingular elliptic curves over  $\mathbb{F}_p$ .
- ▶ So finding a path between two supersingular elliptic curves over  $\mathbb{F}_p$  should be easier than the general problem.
- ▶ Can reduce general case to this case using random walks.
- ▶ We solve the sub-problem using CM theory and algorithm from S. Galbraith, F. Hess, N. P. Smart, “Extending the GHS Weil descent attack”, EUROCRYPT 2002.

# Full supersingular isogeny graph



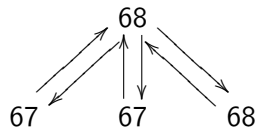
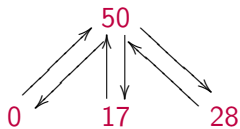
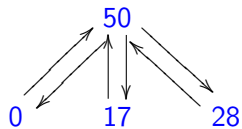
Supersingular Isogeny Graph  $X(\overline{\mathbb{F}}_{83}, 2)$

# Subgraph



Subgraph consisting  $j \in \mathbb{F}_{83}$

# New graph

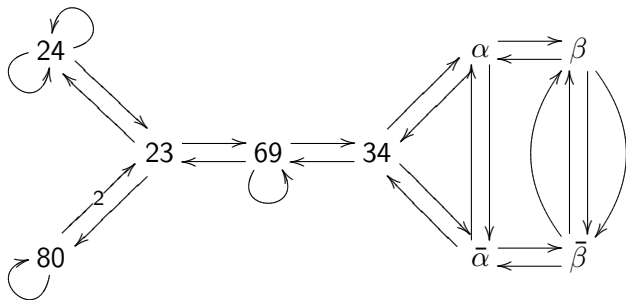


$$X(\mathbb{F}_{83}, 2)$$

## Structure theorem ( $p > 3$ prime)

1.  $p \equiv 1 \pmod{4}$ : There are  $h(-4p)$   $\mathbb{F}_p$ -isomorphism classes of supersingular elliptic curves over  $\mathbb{F}_p$ , all having the same endomorphism ring  $\mathbb{Z}[\sqrt{-p}]$ . From every one there is one outgoing  $\mathbb{F}_p$ -rational horizontal 2-isogeny as well as two horizontal  $\ell$ -isogenies for every prime  $\ell > 2$  with  $\left(\frac{-p}{\ell}\right) = 1$ .
2.  $p \equiv 3 \pmod{4}$ : There are two levels in the supersingular isogeny graph. From each vertex there are two horizontal  $\ell$ -isogenies for every prime  $\ell > 2$  with  $\left(\frac{-p}{\ell}\right) = 1$ .
  - 2.1 If  $p \equiv 7 \pmod{8}$ , on each level  $h(-p)$  vertices are situated. Surface and floor are connected 1:1 with 2-isogenies and on the surface we also have two horizontal 2-isogenies from each vertex.
  - 2.2 If  $p \equiv 3 \pmod{8}$ , we have  $h(-p)$  vertices on the surface and  $3h(-p)$  on the floor. Surface and floor are connected 1:3 with 2-isogenies, and there are no horizontal 2-isogenies.

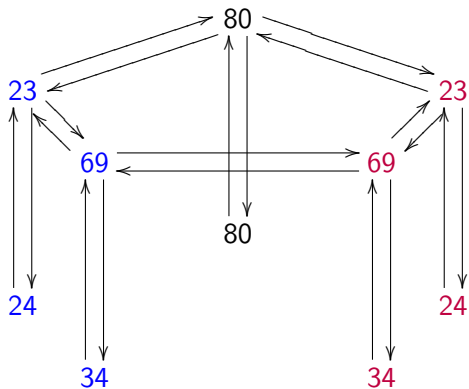
Example 2:  $p = 103 \equiv 7 \pmod{8}$



Supersingular Isogeny Graph  $X(\overline{\mathbb{F}}_{103}, 2)$



New



$X(\mathbb{F}_{103}, 2)$

Thank You

