Completely reducible subgroups of $GL(d, p^{f})$: counting composition factors of order p(joint work with M. Giudici, C. H. Li and G. Verret)

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Outline

- 1 Motivation
- 2 Examples show bounds are best possible
- **3** Aschbacher's classification
- **4** Proof of the main theorem (Thm 4)
- **6** Concluding remarks







Motivation

• Permutation group $G \leq \text{Sym}(\Omega) \rightsquigarrow \text{digraph } \Gamma$.



- (α, β) ∈ Ω × Ω; Arcs of Γ = (α, β)^G → Γ arc transitive.
- $(\beta, \gamma) \in (\alpha, \beta)^{\mathcal{G}}$; $A := \ln \mathbb{N}(\beta)$, $C := \operatorname{OutN}(\beta)$; $L := \mathcal{G}_{\beta}^{\mathcal{A}}$, $R := \mathcal{G}_{\beta}^{\mathcal{C}}$.
- **Theorem [Knapp 1973]** If *L* and *R* are q.p., then *R* is an epimorphic image of *L* or conversely.

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- **Theorem [Knapp 1973]** If *L* and *R* are q.p., then *R* is an epimorphic image of *L* or conversely.
- Suppose R = L/N. There are 8^2 possible types for the pair (L, R) of q.p. groups. It turns out that very few possibilities arise. To eliminate the (funny) possibility (HA, HA) it seemed desirable to prove:
- Theorem [us] If G ≤ GL(d, p) is irreducible, then the number of composition factors of G of order p is at most d − 1.

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- Definition. If G is a finite group, then let c_p(G) denote the number of composition factors of G that have order p.
- Ex 1. $|GL_2(3)| = 2^4 \cdot 3 \quad \rightsquigarrow \quad c_2(GL_2(3)) = 4 \quad c_3(GL_2(3)) = 1.$
- Ex 2. $c_p(T) = 0$ for T nonabelian simple.
- Ex 3. $c_p(GL(d, p^f)) = 0$ if $(d, p^f) \neq (2, 2)$ or (2, 3).
- Ex 4. $c_p(G) \leq \log_p |GL(d, p^f)|_p = {d \choose 2} f$ bounded by the size of Sylow *p*-subgroup.
- Want. If $G \leq GL(d, p)$ is irreducible, then $c_p(G) \leq d 1$.

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- Want. If $G \leq GL(d, p)$ is irreducible, then $c_p(G) \leq d 1$.
- Thm 1. If $G \leq GL(d, p^f)$ is c.r., then $c_p(G) \leq (d-1)f$.
- Thm 2. If $G \leq \operatorname{GL}(d, p^f)$ is c.r., then $c_p(G) \leq (d-1)f/(p-1)$.
- Thm 3. If $G \leq \operatorname{GL}(d, p^f)$ is c.r., then $c_p(G) \leq (\frac{3d}{2} 1)/(p 1)$.
- Thm 4. If $G \leqslant \operatorname{GL}(d, p^f)$ is c.r., then $c_p(G) \leqslant (\varepsilon_{p^f}d 1)/(p 1)$

here
$$\varepsilon_{p^f} = \begin{cases} 4/3 & \text{if } p = 2 \text{ and } f \text{ is even,} \\ p/(p-1) & \text{if } p \text{ is a Fermat prime,} \\ 1 & \text{otherwise.} \end{cases}$$

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Examples show bounds are best possible

- Examples \rightsquigarrow bounds are tight infinitely often.
- Fix $\Gamma_1 \leq \operatorname{GL}(k, p^f)$ and form imprimitive wreath products $\Gamma_n := \Gamma_1 \wr \operatorname{C}_p \wr \cdots \wr \operatorname{C}_p \leq \operatorname{GL}(kp^{n-1}, p^f)$ with n-1 copies of C_p .

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- Generic $\varepsilon_q \ge 1$. Let $q = p^f$ and $\Gamma_1 = C_{p^p-1} \rtimes C_p \le GL(p, p)$, so k = p. Then $\Gamma_n \le GL(p^n, p) \le GL(p^n, q)$ and $c_p(\Gamma_n) = (p^n 1)/(p 1) = (d 1)/(p 1)$.
- $\varepsilon_q \ge p/(p-1)$. If $p = q = 2^m + 1$ is a Fermat prime and Γ_1 is Sylow *p*-subgroup of $GO^-(2m, 2) \le GL(2^m, p) = GL(p-1, p)$, then $\Gamma_n \le GL(d, p)$ is irreducible and $c_p(G) = (p^n - 1)/(p-1)$, so $c_p(\Gamma_n) = (\varepsilon d_n - 1)/(p-1)$ where $d_n = (p-1)p^{n-1}$ and $\varepsilon = p/(p-1)$.
- $\varepsilon_q \ge 4/3$. Take p = 2, $q = 2^2$, and $\Gamma_1 = GU(3, 2)$. Then $\Gamma_n \le GL(3 \cdot 2^{n-1}, 4)$ is irreducible and $c_2(G) = 2^{n+1} - 1$, so $c_p(G) = (\varepsilon d_n - 1)/(p - 1)$ where $d_n = 3 \cdot 2^{n-1}$ and $\varepsilon = 4/3$.

Aschbacher's classification

Dynkin-Aschbacher Theorem. Every completely reducible subgroup G of GL(d, q) lies in at least on of the following classes.

- C_1 (reducible subgps) $V = V_1 \oplus V_2$, $G \leq GL(V_1) \times GL(V_2)$.
- C_2 (imprimitive subgps) $V = V_1 \oplus \cdots \oplus V_r$, $G \leq GL(d/r, q) \wr Sym(r)$.
- \mathcal{C}_3 (ext field subgps) $V = (\mathbb{F}_{q^r})^{d/r}$, and $G \leqslant \operatorname{GL}(d/r, q^r) \rtimes \mathsf{C}_r$.
- C_4 (tensor reducible subgps) $V = V_1 \otimes V_2$ and $G \leq GL(V_1) \otimes GL(V_2)$.
- \mathcal{C}_5 (proper subfield subgps) $G \leqslant \mathsf{GL}(d,q_0) \circ \mathsf{Z}(\mathsf{GL}(d,q)), \ q = q_0^r$.
- C_6 (symplectic type *r*-groups) $d = r^m$, $G \leq N_{GL(d,q)}(R)$ where $R/Z(R) \cong C_r^{2m}$ is elementary, and $\Phi(R) \leq Z(R)$.
- C_7 (tensor reducible subgps) $V = V_1 \otimes \cdots \otimes V_r$ and $G \leq GL(V_1) \wr Sym(r)$.
- C₈ (classical groups) preserves symplectic, unitary, or orthogonal form and contains Sp(V)', SU(V), or Ω^ε(V) resp., where ε ∈ {±, ∘}.
- C₉ (nearly simple) Z := Z(G), socle(G/Z) = N/Z is almost simple and absolutely irreducible.

Proof of the main theorem (Thm 4)

• Induction on (d, q) ordered lexicographically

 $(d_1, q_1) < (d_2, q_2)$ if $d_1 < d_2$ or $d_1 = d_2$ and $q_1 < q_2$.

- Simple cases:
- \mathcal{C}_1 . Then $G \leqslant \operatorname{GL}(d_1,q) \times \operatorname{GL}(d_2,q)$, so $G \leqslant G_1 \times G_2$ and

$$c_{p}(G) \leq c_{p}(G_{1}) + c_{p}(G_{2}) \leq \frac{\varepsilon_{q}d_{1} - 1}{p - 1} + \frac{\varepsilon_{q}d_{1} - 1}{p - 1}$$
$$= \frac{\varepsilon_{q}(d_{1} + d_{2}) - 2}{p - 1} < \frac{\varepsilon_{q}d - 1}{p - 1}.$$

• C_4 . Then $G \leq \operatorname{GL}(d/r,q) \wr \operatorname{Sym}(r)$, so $G \leq G_1 \wr G_2$ and

$$c_{\rho}(G) \leqslant rc_{\rho}(G_1) + c_{\rho}(G_2) \leqslant rac{r(arepsilon_q d/r - 1)}{p-1} + rac{r-1}{p-1} = rac{arepsilon_q d-1}{p-1}$$

Proof of the main theorem (Thm 4)

- C_2 . Like C_1 ; C_3 and C_5 . Induction; C_7 . Like C_4 .
- C_8 . Simple, literally.
- Harder case:
- \mathcal{C}_6 . Number theory |G| small $\rightsquigarrow c_p(G)$ small.
- Hardest case:
- C_9 . T = N/Z simple, |G/N| divides |Out(T)|, $c_p(G) = \log_p |G/N|_p \leq \log_p |Out(T)|_p$. Most difficulties when T = L(q') simple of Lie-type.

Concluding remarks

- Obtain insight into local symmetries of digraphs.
- Apply results to limit the local symmetries of digraphs, and construct new highly symmetric examples.
- What if the prime p ≠ char(F_q)? If G ≤ GL(d, q) is completely reducible, then find sharp upper bounds for c_p(G). (Partially solved.)
- Counting other composition factor. (Partial progress.)

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Thank You!