A Counterexample to Fishburn's Conjecture

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Abstract

Kraft, Pratt and Seidenberg (1959) provided an infinite set of axioms which, when taken together with de Finetti's axiom, gives a necessary and sufficient set of "cancellation" conditions for representability of an ordering relation on subsets of a set by an order-preserving probability measure. Fishburn (1996) defined $f(n)$ to be the smallest positive integer k such that every comparative probability ordering on an n-element set which satisfies the cancellation conditions $C_4, \dots, C_k$ is representable. By the work of Kraft, Pratt and Seidenberg (1959) and Fishburn (1996, 1997), it is known that $n - 1 \leq f(n) \leq n + 1$ for all $n \geq 5$. Also Fishburn proved that $f(5) = 4$, and conjectured that $f(n) = n - 1$ for all $n \geq 5$. In this paper we confirm that $f(6) = 5$, but give counter-examples to Fishburn's conjecture for $n = 7$, showing that $f(7) \geq 7$. We summarise, correct and extend many of the known results on this topic, including the notion of "almost representability", and offer an amended version of Fishburn's conjecture.

1 Introduction

The topic of this article dates back to fundamental work by Bruno de Finetti [3], who asked whether a certain axiom for a comparative probability ordering relation $\succeq$ on the subsets of a set is sufficient for the existence of an order-preserving probability measure on that set. For infinite sets, assuming there is a uniform partition into an arbitrary number of events, Savage [10] answered this question positively. For finite sets of cardinality 5 or more, a negative answer to this question was given by Kraft, Pratt and Seidenberg [6]. It became clear from these observations that comparative probability is a broader concept than probability.

Nevertheless, de Finetti's axioms were very attractive, especially for use in applications where relative frequency characterization of probabilities is not available or meaningful, for example when we have to deal with the so-called "subjective probability" (also called "intuitive probability" or "personal probability"), which is essentially a belief of an individual that some events are more probable than the others [8]. This led to development of comparative probability as a separate area of research (see a recent survey of Giuliana Regoli [9]).

Sometimes comparative probability orderings emerge in circumstances where we cannot speak even about subjective probability. For example, they appeared in an article by Vladimir Danilov [2], who, in the case of dichotomous preferences, characterised all social choice rules which satisfy the famous Arrow's condition of Independence of Irrelevant Alternatives. It appeared that each such rule corresponds to a comparative probability ordering, and the rule is a weighted majority rule if and only if the comparative probability ordering is representable by an order-preserving probability measure on the set of voters.

Kraft, Pratt and Seidenberg [6] discovered an infinite set of axioms which, when taken together with de Finetti's axiom, gives a necessary and sufficient set of conditions for representability of an ordering relation on subsets by an order-preserving probability measure. Implicitly, these axioms involved a concept of a multiset, which was not widely understood at the time, and reformulation of these axioms by Scott [11] helped to clarify things. Krantz, Luce, Suppes and Tversky [7] called these axioms cancellation conditions. Peter Fishburn [4, 5] opined that they "play a central role in the representational theory of measurement", and devoted much effort to their study.

Fishburn introduced a function $f(n)$, which for each positive integer n measures the maximal deviation of a comparative probability ordering on an n-element set from one which arises from an order-preserving probability measure. It follows from [6] that $f(n) \leq n+1$. Fishburn then proved in [4, 5]

that $f(n) \geq n - 1$ for all $n \geq 5$, and so combining these two results gives the following bounds for this function: $n - 1 \leq f(n) \leq n + 1$. Fishburn also proved that $f(5) = 4$, and conjectured that $f(n) = n - 1$ for all $n \geq 5$.

In this paper we confirm that $f(6) = 5$, but give counter-examples to Fishburn's conjecture for $n = 7$ showing that $f(7) \geq 7$. We summarise all known results, correcting and extending some of them, and offer an amended version of Fishburn's conjecture.

2 Preliminaries

To formulate cancellation conditions, let us first recall some other definitions and fix some notation. Here 2^X denotes the set of all subsets of a set X.

Definition 1. *Let S be a finite set. Any reflexive, complete and transitive binary relation $\succeq$ on S will be called an order on S. (Note: sometimes this is called a complete pre-order.) Such an order $\succeq$ gives rise to two other relations $\succ$ and $\sim$ on S, defined for all $x, y \in S$ by*

(a) $x \succ y \iff x \succeq y$ *and* $\text{not}\,(y \succeq x)$;

(b) $x \sim y \iff x \succeq y$ *and* $y \succeq x$.

If an order $\succeq$ is anti-symmetric, then all equivalence classes (under the relation $\sim$) are singletons, and $\succeq$ is called a linear order.

Definition 2. *Let X be a finite set. A linear order $\succeq$ on 2^X is called a comparative probability ordering on X, if $A \succ \emptyset$ for every non-empty subset A of X, and $\succeq$ satisfies de Finetti's axiom, namely for all $A, B, C \in 2^X$,*

$$A \succeq B \iff A \cup C \succeq B \cup C \text{ whenever } (A \cup B) \cap C = \emptyset. \qquad (1)$$

For convenience, we will further suppose that $X = I_n = \{1, 2, \ldots, n\}$. Note that if we have a probability measure $\mathbf{p} = (p_1, \ldots, p_n)$ on X, where p_i is the probability of i, then we know the probability of every event A, by the rule $p(A) = \sum_{i \in A} p_i$. We may now define a relation $\succeq$ on 2^X by

$$A \succeq B \quad \text{if and only if} \quad p(A) \geq p(B).$$

If $p_i > 0$ for all i, and the probabilities of all events are different, then $\succeq$ is a comparative probability ordering on X, and any such ordering is called *representable* [9]. On the other hand, any comparative probability ordering $\succeq$ which is not obtainable in this way is called *non-representable*.

It can also be possible that for some measure $\mathbf{p}$ on X the probabilities of two or more events coincide, that is $p(A) = p(B)$ when $A \succeq B$ but $A \neq B$.

In this case we have an ordering $\succeq$ which still satisfies the de Finetti axiom (1), and we call it a *weak comparative probability ordering* arising from **p**.

If we have a non-representable comparative probability ordering $\succeq$ on X, then we cannot specify the exact probability of every event, but if $A \succ B$ then we can say that A is more probable than B. The class of comparative probability measures is therefore broader than the class of probability measures [6]. A non-representable comparative probability ordering $\succeq$ on X is said to *almost agree* with the measure **p** on X if

$$A \succeq B \Longrightarrow p(A) \geq p(B). \tag{2}$$

If such a measure **p** exists, then the ordering $\succeq$ is said to be *almost representable*. When **p** defines a weak comparative probability ordering on X, the ordering $\succeq$ may not be representable.

We now recall that a set is a special case of a multiset. Multisets are collections of objects which may include several copies of the same object. If we consider multisets on X, then we allow an element $x \in X$ to enter such a multiset M with multiplicity $\mu_M(x)$, which is a non-negative integer. Multiset union $\cup$ is an operation over multisets that adds multiplicities of elements, that is, $\mu_{A \cup B}(x) = \mu_A(x) + \mu_B(x)$. See Stanley [12] for more information on multisets.

Definition 3. *A linear order $\succeq$ is said to satisfy the mth cancellation condition C_m if for no m distinct comparisons $A_i \succ B_i$ (for $1 \leq i \leq m$), there exist positive integers $a_1, \ldots, a_m$ such that the following two multiset unions coincide:*

$$\bigcup_{i=1}^{m} \underbrace{(A_i \cup \ldots \cup A_i)}_{a_i} = \bigcup_{i=1}^{m} \underbrace{(B_i \cup \ldots \cup B_i)}_{a_i}. \tag{3}$$

Alternatively, in the spirit of [11] (and also [13]), condition (3) can be reformulated in terms of characteristic functions of subsets, as follows:

$$\sum_{i=1}^{m} a_i \chi_{A_i} = \sum_{i=1}^{m} a_i \chi_{B_i},$$

where the characteristic function χ_S of the set S is given by $\chi_S(x) = 1$ if $x \in S$ and $\chi_S(x) = 0$ if $x \notin S$.

Example 1. *The non-representable linear order $\succeq$ on I_5 constructed in [6] does not satisfy the condition C_4 since it contains the following comparisons:*

$$\{2, 4, 5\} \prec \{1, 3\}, \ \{1, 5\} \prec \{2, 4\}, \ \{3, 4\} \prec \{2, 5\}, \ \{2\} \prec \{4, 5\}. \tag{4}$$

Indeed, the multiset union of the sets on the right and the multiset union of the sets on the left are both equal to the multiset $\{1, 2^2, 3, 4^2, 5^2\}$, where powers denote multiplicities. Hence the cancellation condition C_4 is violated with $a_1 = a_2 = a_3 = a_4 = 1$.

Proposition 1 (Kraft-Pratt-Seidenberg [6]). *For a comparative probability ordering to be representable, it is necessary and sufficient that all cancellation conditions C_m are satisfied.*

These cancellation conditions look rather unnatural and complicated. Since they are derived from linear algebra, it is not surprising that they can be made to look much more natural by being reformulated in terms of vectors, as we do later in this paper.

It is clear that any representable linear order satisfies C_m for all $m \geq 1$. It was also shown in [4] that C_2 and C_3 follow from de Finetti's axiom and properties of linear orders. Hence C_4 is the first nontrivial cancellation condition.

As was noticed in [6], for $n < 5$ all comparative probability orderings are representable, but for $n = 5$ there are non-representable ones. Fishburn [4] showed that all non-representable comparative probability orderings on a 5-element set X fail to satisfy C_4, and conjectured [4, 5] that any such ordering on an n-element set X, which satisfies $C_4, \ldots, C_{n-1}$, is representable.

In Section 3 we will exhibit counter-examples to this conjecture for $n = 7$. More precisely, we will construct non-representable comparative probability orderings on a 7-element set X which satisfy C_4, C_5 and C_6. In addition, we will show these orderings are almost representable.

We will always assume here that a linear order $\succeq$ on 2^X is a comparative probability ordering. As in [4, 5], we will also assume that $1 \succ 2 \succ \ldots \succ n$, which is equivalent to assuming that $p_1 > p_2 > \ldots > p_n$ for a comparative probability ordering arising from a probability measure $\mathbf{p} = (p_1, \ldots, p_n)$. To every such linear order $\succeq$, there corresponds a *discrete cone* $C(\succeq)$ in T^n, where $T = \{-1, 0, 1\}$ (as defined in [4]).

Definition 4. *A subset $C \subseteq T^n$ is said to be a discrete cone if the following properties hold:*

D1. $\{\mathbf{e}_1 - \mathbf{e}_2, \ldots, \mathbf{e}_{n-1} - \mathbf{e}_n, \mathbf{e}_n\} \subseteq C$, *where $\{\mathbf{e}_1, \ldots, \mathbf{e}_n\}$ is the standard basis of $\mathbb{R}^n$,*

D2. $\{-\mathbf{x}, \mathbf{x}\} \cap C \neq \emptyset$ *for every $\mathbf{x} \in T^n$,*

D3. $\mathbf{x} + \mathbf{y} \in C$ *whenever $\mathbf{x}, \mathbf{y} \in C$ and $\mathbf{x} + \mathbf{y} \in T^n$.*

We note that in [4], Fishburn requires $\mathbf{0} \notin C$ because his orders are anti-reflexive. In our case, condition D2 implies $\mathbf{0} \in C$.

Given a (weak) comparative probability ordering $\succeq$ on X, we may construct a characteristic vector $\chi(A, B) = \chi_A - \chi_B \in T^n$ for every possible comparison $A \succeq B$. The cone $C(\succeq)$ is then defined as the set of all characteristic vectors $\chi(A, B)$, for $A, B \in 2^X$ such that $A \succeq B$. The two axioms of comparative probability guarantee that $C(\succeq)$ is indeed a discrete cone (see [4, Lemma 2.1]).

Example 2. *For the comparative probability ordering $\succeq$ from Example 1, the four vectors*

$$(1, -1, 1, -1, -1), \ (-1, 1, 0, 1, -1), \ (0, 1, -1, -1, 1), \ (0, -1, 0, 1, 1) \quad (5)$$

all belong to $C(\succeq)$, and correspond to the four comparisons listed in (4) above.

We will say that a discrete cone $C \subseteq T^n$ is *generated* by a set of vectors $V = \{\mathbf{v}_1, \ldots, \mathbf{v}_m\} \subseteq T^n$ if C is the smallest discrete cone containing V. It follows that the cone $C(\succeq)$ associated with a comparative probability ordering $\succeq$ must always contains the discrete cone U in T^n generated by

$$\{\mathbf{e}_1 - \mathbf{e}_2, \ \ldots, \mathbf{e}_{n-1} - \mathbf{e}_n, \ \mathbf{e}_n\}. \quad (6)$$

We observe that a vector $\mathbf{x} = (x_1, \ldots, x_n) \in T^n$ belongs to U if and only if

$$x_1 \geq 0, \ x_1 + x_2 \geq 0, \ldots, x_1 + x_2 + \ldots + x_n \geq 0. \quad (7)$$

Hence we can reformulate the cancellation conditions as follows [4]:

Proposition 2. *A comparative probability ordering $\succeq$ satisfies the mth cancellation condition C_m if and only if for no set $\{\mathbf{x}_1, \ldots, \mathbf{x}_m\}$ of non-zero vectors in $C(\succeq)$ do there exist positive integers $a_1, \ldots, a_m$ such that*

$$a_1 \mathbf{x}_1 + a_2 \mathbf{x}_2 + \cdots + a_m \mathbf{x}_m = 0. \quad (8)$$

Proof. Take a vector $\mathbf{x}_i = \chi(A_i, B_i) \in C(\succeq)$ for each comparison $A_i \succeq B_i$ (for $1 \leq i \leq m$) involved in Definition 3, and observe that condition (3) is equivalent to the equation $a_1 \mathbf{x}_1 + a_2 \mathbf{x}_2 + \cdots + a_m \mathbf{x}_m = \mathbf{0}$. $\square$

Example 3. *The four vectors from the cone $C(\succeq)$ given in Example 2 add up to the zero vector, hence the corresponding ordering $\succeq$ is not representable.*

Geometrically, what happens is clear. A comparative probability ordering $\succeq$ is representable if and only if there exists a positive integer-valued vector $\mathbf{a} \in \mathbb{R}^n$ such that

$$\mathbf{x} \in C(\succeq) \Longleftrightarrow (\mathbf{a}, \mathbf{x}) > 0 \quad \text{for all } \mathbf{x} \in T^n \setminus \{\mathbf{0}\},$$

that is, every vector in the cone $C(\succeq)$, apart from the zero vector, lies in the open half-space $S_{\mathbf{a}} = \{\mathbf{x} \in \mathbb{R}^n \mid (\mathbf{a}, \mathbf{x}) > 0\}$.

Similarly, for a non-representable but *almost* representable comparative probability ordering $\succeq$, there exists an integer-valued vector $\mathbf{a} \in \mathbb{R}^n$ with non-negative entries such that

$$\mathbf{x} \in C(\succeq) \Longleftrightarrow (\mathbf{a}, \mathbf{x}) \geq 0 \quad \text{for all } \mathbf{x} \in T^n \setminus \{\mathbf{0}\},$$

and hence in this case the whole cone $C(\succeq)$ must lie in the closed half-space $\overline{S_{\mathbf{a}}} = \{\mathbf{x} \in \mathbb{R}^n \mid (\mathbf{a}, \mathbf{x}) \geq 0\}$.

In both cases, the normalised vector $\mathbf{a}$ gives us the probability measure, namely $\mathbf{p} = (a_1 + \ldots + a_n)^{-1} (a_1, \ldots, a_n)$, from which $\succeq$ arises or with which it almost agrees.

3   Almost representable orderings

Proposition 3. *Let $\succeq$ be a non-representable but almost representable comparative probability ordering which almost agrees with a probability measure $\mathbf{p}$. Suppose that the mth cancellation condition C_m is violated, and that for some non-zero vectors $\{\mathbf{x}_1, \ldots, \mathbf{x}_m\} \subseteq C(\succeq)$ and some positive integers $a_1, \ldots, a_m$, the condition (8) holds. Then all of the vectors $\mathbf{x}_1, \ldots, \mathbf{x}_m$ lie in the hyperplane $H_{\mathbf{p}} = \{\mathbf{x} \in \mathbb{R}^n \mid (\mathbf{p}, \mathbf{x}) = 0\}$.*

Proof. First note that for every $\mathbf{x} \in C(\succeq)$ which does not belong to $H_{\mathbf{p}}$, we have $(\mathbf{p}, \mathbf{x}) > 0$. Hence the condition (8) can hold only when all $\mathbf{x}_i \in H_{\mathbf{p}}$. $\square$

Corollary 1. *Any almost representable comparative probability ordering $\succeq$ on an n-element set X, which satisfies $C_4, \ldots, C_n$, is representable.*

Proof. Suppose that (8) holds for some non-zero vectors $\mathbf{x}_1, \ldots, \mathbf{x}_m \in C(\succeq)$ and positive integers $a_1, \ldots, a_m$, where $m > 0$. As all m vectors $\mathbf{x}_1, \ldots, \mathbf{x}_m$ lie in the $(n-1)$-dimensional subspace $H_{\mathbf{p}}$, we may, using standard linear algebra, reduce the number of vectors in this linear combination to at most n vectors while still having positive coefficients. Hence we may suppose that $m \leq n$. However, this is impossible since we assumed that C_m is satisfied. Therefore $\succeq$ must be representable. $\square$

Following Fishburn [4], we now define $f(n)$ to be the smallest positive integer k such that every comparative probability ordering on an n-element set which satisfies the cancellation conditions $C_4, \ldots , C_k$ is representable. Fishburn proved in [4] that $f(5) = 4$, and in [5] that $n - 1 \leq f(n) \leq n + 1$ for all $n \geq 6$, and conjectured that $f(n) = n - 1$ for all $n \geq 5$ (see [5, p.354]).

In the spirit of Kraft, Pratt and Seidenberg [6], we further define $g(n)$ to be the smallest positive integer k such that every *almost resentable* comparative probability ordering on an n-element set which satisfies the cancellation conditions $C_4, \ldots , C_k$ is representable. Clearly $g(n) \leq f(n)$ for all n, and also by Corollary 1 above, we have the following:

Corollary 2. $g(n) \leq n$ for all n.

Now Fishburn's conjecture, if valid, would imply that $g(n) \leq n - 1$ for all n. We can confirm that $f(6) = 5$, which is consistent with Fishburn's conjecture, but we also prove that $g(7) = 7$, which refutes it. To do this, we will need the following theorem which simplifies and generalises Theorem 6.1 from [4] (and its generalisation to Lemma 1 in [5]).

Theorem 1 (Construction method 1). *Let $X = \{\mathbf{x}_1, \ldots , \mathbf{x}_m\}$, $m \geq 4$, be a system of non-zero vectors from T^n, such that $\sum_{i=1}^{m} a_i \mathbf{x}_i = 0$ for some positive integers $a_1, \ldots , a_m$, and such that no proper subsystem $X' \subset X$ is linearly dependent with positive coefficients. Suppose further that the $m \times n$ matrix A having the vectors $\mathbf{x}_1, \ldots , \mathbf{x}_m$ as its rows has the property that $A\mathbf{b} = \mathbf{0}$ for some positive integer-valued vector $\mathbf{b} = (b_1, \ldots , b_n)$ with $b_1 > b_2 > \ldots > b_n > 0$, and that*

$$Span\{\mathbf{x}_1, \ldots , \mathbf{x}_m\} \cap T^n = \{\pm\mathbf{x}_1, \ldots , \pm\mathbf{x}_m\}.$$

Let $C(\succeq)$ be the discrete cone belonging to the weak comparative probability ordering which arises from the measure $\mathbf{p} = (b_1 + \ldots + b_n)^{-1} (b_1, \ldots , b_n)$, that is, $C(\succeq) = \{\mathbf{x} \in \mathbb{R}^n \mid (\mathbf{x}, \mathbf{b}) \geq 0\}$. Then the discrete cone

$$C' = C(\succeq) \setminus \{-\mathbf{x}_1, \ldots , -\mathbf{x}_m\}$$

corresponds to an almost representable comparative probability ordering which almost agrees with $\mathbf{p}$, and satisfies C_i for all $i < m$, but does not satisfy C_m.

Proof. We first note that $\{\pm\mathbf{x}_1, \ldots , \pm\mathbf{x}_m\} \subseteq H_\mathbf{b} = \{\mathbf{x} \in \mathbb{R}^n \mid (\mathbf{b}, \mathbf{x}) = 0\}$. Next $(\mathbf{e}_i - \mathbf{e}_{i+1}, \mathbf{b}) = b_i - b_{i+1} > 0$ for $1 \leq i < n$, and also $(\mathbf{e}_n, \mathbf{b}) = b_n > 0$, so the property D1 holds for the cone C'. Property D2 holds for C' because it holds for $C(\succeq)$ and when we remove vectors $-\mathbf{x}_1, \ldots , -\mathbf{x}_m$ from C' the vectors $\mathbf{x}_1, \ldots , \mathbf{x}_m$ remain in C'. Finally let us prove D3. Suppose that

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$\mathbf{y}, \mathbf{z} \in C'$ and $\mathbf{y} + \mathbf{z} \in T^n$, but $\mathbf{y} + \mathbf{z} \notin C'$. Since $\mathbf{y} + \mathbf{z} \in C(\succeq) \setminus C'$, we deduce that $\mathbf{y} + \mathbf{z} = -\mathbf{x}_i$ for some i. Moreover, $\mathbf{y}, \mathbf{z} \in H_\mathbf{b}$, since otherwise $(\mathbf{y} + \mathbf{z}, \mathbf{b}) = (\mathbf{y}, \mathbf{b}) + (\mathbf{z}, \mathbf{b}) > 0$. Hence $\mathbf{y} = \mathbf{x}_j$ for some j and $\mathbf{z} = \mathbf{x}_k$ for some k, but from this it follows that $\mathbf{x}_i + \mathbf{x}_j + \mathbf{x}_k = \mathbf{0}$, a contradiction since $|X| = m \geq 4$ and hence no three rows of A may have sum zero.

Suppose now that some violation of C_k occurs, say

$$c_1 \mathbf{y}_1 + c_2 \mathbf{y}_2 + \cdots + c_k \mathbf{y}_k = \mathbf{0},$$

with positive integers $c_1, \ldots, c_k$ and $\mathbf{y}_i \in C(\succeq)$ for $1 \leq i \leq k$, where $k < m$. If $(\mathbf{y}_i, \mathbf{b}) > 0$ from some i, then $(\mathbf{0}, \mathbf{b}) = (c_1 \mathbf{y}_1 + c_2 \mathbf{y}_2 + \cdots + c_k \mathbf{y}_k, \mathbf{b}) = c_1(\mathbf{y}_1, \mathbf{b}) + c_2(\mathbf{y}_2, \mathbf{b}) + \ldots + c_k(\mathbf{y}_k, \mathbf{b}) > 0$, a contradiction, hence $\mathbf{y}_i \in H_\mathbf{b}$ for all i, and therefore $\{\mathbf{y}_1, \ldots, \mathbf{y}_k\} \subseteq \{\mathbf{x}_1, \ldots, \mathbf{x}_m\}$, which is impossible. On the other hand, C_m fails because $\sum_{i=1}^{m} a_i \mathbf{x}_i = 0$. $\qquad\square$

A major difference between our theorem and Theorem 6.1 of [4] is that we claim almost representability of the constructed ordering. Using this, we can now prove the following:

Theorem 2. $f(7) \geq g(7) = 7.$

Proof. Consider the following 7×7 matrix

$$A = \begin{bmatrix} -1 & 1 & 1 & -1 & 0 & 0 & 0 \\ 1 & 0 & -1 & -1 & 1 & 0 & 0 \\ 1 & -1 & 1 & 0 & 0 & -1 & -1 \\ -1 & -1 & 0 & 1 & 1 & 0 & 1 \\ -1 & 1 & -1 & 1 & 0 & 1 & -1 \\ 1 & -1 & 1 & -1 & -1 & 1 & 0 \\ 0 & 1 & -1 & 1 & -1 & -1 & 1 \end{bmatrix}.$$

Denote the rows of A by $\mathbf{x}_1, \ldots, \mathbf{x}_7$. As the null space of A^T is spanned by the vector $(1, 1, 1, 1, 1, 1, 1)$, we can see that $\mathbf{x}_1 + \mathbf{x}_2 + \ldots + \mathbf{x}_7 = \mathbf{0}$ and that no six rows of A are linearly dependent. We also observe that $A\mathbf{b} = \mathbf{0}$ for $\mathbf{b} = (27, 25, 24, 22, 19, 15, 11)$. Let us take the probability measure $\mathbf{p} = \frac{1}{116}(27, 25, 24, 22, 19, 15, 11)$, and now consider the weak comparative probability ordering on $I_7 = \{1, 2, 3, 4, 5, 6, 7\}$ that arises from $\mathbf{p}$.

Ignoring braces when writing subsets, and listing only half of all terms (as $A \succeq B$ if and only if $\bar{B} \succeq \bar{A}$), we may describe this ordering by the following sequence (given in reverse order for simplicity):

$\emptyset \prec 7 \prec 6 \prec 5 \prec 4 \prec 3 \prec 2 \prec 67 \prec 1 \prec 57 \prec 47 \prec 56 \prec 37 \prec 27 \prec 46 \prec$
$17 \prec 36 \prec 26 \prec 45 \prec 16 \prec 35 \prec 25 \prec 567 \prec 15 \sim 34 \prec 24 \prec 467 \prec 23 \sim$
$14 \prec 367 \prec 13 \sim 267 \prec 457 \sim 12 \prec 167 \prec 357 \prec 257 \prec 456 \prec 157 \sim 347 \prec$

$247 \sim 356 \prec 256 \prec 237 \sim 147 \prec 156 \sim 346 \prec 246 \sim 137 \prec 127 \prec 236 \sim$
$146 \prec 345 \prec 136 \sim 245 \prec 4567 \sim 126 \prec 235 \sim 145 \prec 3567 \prec 135 \sim 2567 \prec$
$125 \sim 234 < \ldots$.

The only equivalences in the first half of this sequence are these:

$$15 \sim 34, \quad 157 \sim 347, \quad 156 \sim 346, \quad 125 \sim 234,$$
$$23 \sim 14, \quad 237 \sim 147, \quad 236 \sim 146, \quad 235 \sim 145,$$
$$13 \sim 267, \quad 135 \sim 2567,$$
$$457 \sim 12, \quad 4567 \sim 126,$$
$$246 \sim 137,$$
$$136 \sim 245,$$
$$247 \sim 356.$$

Note that in each row of the above list, all equivalences are the consequences of the leftmost equivalence, and all contribute the same pair of vectors $\pm \mathbf{x}_i$ to the cone $C(\succeq)$, where $\mathbf{x}_1, \ldots, \mathbf{x}_7$ are the rows of A. All equivalences from the second half of the sequence are also consequences of these. It follows that

$$Span\{\mathbf{x}_1, \ldots, \mathbf{x}_7\} \cap T^7 = \{\pm \mathbf{x}_1, \ldots, \pm \mathbf{x}_7\}.$$

We may now apply Theorem 1, to deduce the existence of a non-representable comparative probability ordering on I_7, which satisfies C_4, C_5 and C_6. $\qquad \square$

Note that the example given in the above proof fails the cancellation condition C_7 with $(a_1, \ldots, a_7) = (1, 1, 1, 1, 1, 1, 1)$. There are hundreds of other examples with the same property. Another interesting example will be given in Section 4.

4 Alternative form of cancellation conditions

In the spirit of Fishburn's earlier paper [4], we now introduce an alternative form of cancellation conditions.

Definition 5. *A comparative probability ordering $\succeq$ is said to satisfy the mth cancellation condition $C_m^\dagger$ if for no set $\{\mathbf{x}_1, \ldots, \mathbf{x}_k\}$ of non-zero vectors in $C(\succeq)$ do there exist positive integers $a_1, \ldots, a_k$ such that*

$$a_1 \mathbf{x}_1 + a_2 \mathbf{x}_2 + \cdots + a_k \mathbf{x}_k = 0 \tag{9}$$

and $\sum_{i=1}^k a_i = m$.

Note here that while the condition C_m emphasises the total number of *distinct* comparisons $A_i \succ B_i$ involved, the alternative condition $C_m^\dagger$ is more concerned with the total number of (not necessarily distinct) comparisons actually required for the sum.

It is clear (and was observed by Fishburn [4]) that a comparative probability ordering satisfies C_m for all $m \geq 4$ if and only if it satisfies $C_m^\dagger$ for all $m \geq 4$. It follows that these alternative cancellation conditions can also be used for measuring the degree of deviation of a comparative probability ordering from one that arises from a probability measure.

Hence we may define $f^\dagger(n)$ to be the smallest positive integer k such that every comparative probability ordering on an n-element set, which satisfies the cancellation conditions $C_4^\dagger, \ldots, C_k^\dagger$, is representable. The corresponding analogue of $g(n)$ can also be introduced. We define $g^\dagger(n)$ to be the smallest positive integer k such that every *almost reresentable* comparative probability ordering on an n-element set, which satisfies the cancellation conditions $C_4^\dagger, \ldots, C_k^\dagger$, is representable. Obviously, $f^\dagger(n) \geq g^\dagger(n)$, and also $f^\dagger(n) \geq f(n)$ and $g^\dagger(n) \geq g(n)$ for all n. Also Kraft, Pratt and Seidenberg effectively commented without proof that $f^\dagger(n) \leq (n+1)!$ (see [6, p. 413]).

Despite such a large upper bound, it appears that to date no examples have been found for which $f^\dagger(n) \neq f(n)$ or $g^\dagger(n) \neq g(n)$. We show below that $f^\dagger(7) \geq g^\dagger(7) \geq 8$, which, together with our earlier finding that $g(7) = 7$, proves that $g^\dagger(7) > g(7)$. To do this, we will need a variation of Theorem 1.

Theorem 3 (Construction method 2). *Let $X = \{\mathbf{x}_1, \ldots, \mathbf{x}_k\}$, $k \geq 4$, be a system of non-zero vectors from T^n such that $\sum_{i=1}^{k} a_i \mathbf{x}_i = 0$ for some non-negative integers $a_1, \ldots, a_k$ with $\sum_{i=1}^{k} a_i = m$, and such that for every equation $\sum_{i=1}^{k} a_i' \mathbf{x}_i = 0$ with positive integer coefficients a_i' it happens that $\sum_{i=1}^{k} a_i' \geq m$. Suppose further that the $k \times n$ matrix A having the vectors $\mathbf{x}_1, \ldots, \mathbf{x}_k$ as its rows has the property that $A\mathbf{b} = \mathbf{0}$ for some positive integer-valued vector $\mathbf{b} = (b_1, \ldots, b_n)$ with $b_1 > b_2 > \ldots > b_n > 0$, and that*

$$Span\{\mathbf{x}_1, \ldots, \mathbf{x}_k\} \cap T^n = \{\pm\mathbf{x}_1, \ldots, \pm\mathbf{x}_k\}.$$

Let $C(\succeq)$ be the discrete cone belonging to the weak comparative probability ordering which arises from the measure $\mathbf{p} = (b_1 + \ldots + b_n)^{-1} (b_1, \ldots, b_n)$, that is, $C(\succeq) = \{\mathbf{x} \in \mathbb{R}^n \mid (\mathbf{x}, \mathbf{b}) \geq 0\}$. Then the discrete cone

$$C' = C(\succeq) \setminus \{-\mathbf{x}_1, \ldots, -\mathbf{x}_m\}$$

corresponds to an almost representable comparative probability ordering which almost agrees with $\mathbf{p}$, and satisfies $C_i^\dagger$ for all $i < m$, but does not satisfy $C_m^\dagger$.

Proof. The proof is completely analogous to the proof of Theorem 1. $\qquad\square$

Theorem 4. $f^\dagger(7) \geq g^\dagger(7) \geq 8$.

Proof. Consider the following 7×7 matrix:

$$B = \begin{bmatrix} 1 & -1 & -1 & 1 & 0 & 1 & -1 \\ 1 & 0 & -1 & -1 & 1 & -1 & -1 \\ 1 & 0 & -1 & -1 & -1 & 0 & 1 \\ -1 & 1 & -1 & 1 & 1 & 0 & 1 \\ 0 & -1 & 1 & 1 & 0 & -1 & 1 \\ 0 & -1 & 1 & -1 & 1 & 1 & 1 \\ -1 & 1 & 1 & 0 & -1 & 0 & -1 \end{bmatrix},$$

and let $\mathbf{x}_1, \ldots, \mathbf{x}_7$ denote its rows. It is easy to check that

$$\mathbf{x}_1 + \ldots + \mathbf{x}_6 + 2\mathbf{x}_7 = \mathbf{0},$$

that is, condition C_7 is violated with $(a_1, \ldots, a_7) = (1, 1, 1, 1, 1, 1, 2)$, and there does not exist any "smaller" violation. Calculations also show that

$$Span\{\mathbf{x}_1, \ldots, \mathbf{x}_7\} \cap T^7 = \{\pm\mathbf{x}_1, \ldots, \pm\mathbf{x}_7\},$$

and $A\mathbf{p} = \mathbf{0}$ for the probability measure $\mathbf{p} = \frac{1}{148}(48, 40, 27, 16, 12, 10, 7)$. For the weak comparative probability ordering determined by this measure, the equivalences are all consequences of one of the following seven: $146 \sim 237$, $15 \sim 3467$, $17 \sim 345$, $2457 \sim 13$, $347 \sim 26$, $3567 \sim 24$ and $23 \sim 157$. Again this is a non-representable but almost representable comparative probability ordering on I_7, but is one which satisfies $C_4^\dagger$, $C_5^\dagger$, $C_6^\dagger$, and $C_7^\dagger$, while failing $C_8^\dagger$ with $(a_1, \ldots, a_7) = (1, 1, 1, 1, 1, 1, 2)$. $\qquad\square$

5 Account for n = 5 and n = 6

The results of this section were obtained with the help of the computer algebra system MAGMA [1], using the same techniques of enumeration as described by Fishburn in [4, Section 4], but also checking the results for both representability and almost representability (using techniques of linear programming). Enumeration of the orderings required less than 60 seconds of computer time for the case $n = 5$, and 12 hours for the case $n = 6$, while determining representability (and/or partial representability) took just a little longer.

Theorem 5. *There are* 546 *different comparative probability orderings of the subsets of* $I_5 = \{1, 2, 3, 4, 5\}$, *of which* 516 *are representable, and the remaining* 30 *are almost representable.*

In particular, our computations show that Fishburn's claimed total number of 561 orderings in [4] is incorrect.

Note also that it is quite easy to decide if any given one of these orderings is almost representable. For example, one of the 30 non-representable orderings (the last in the list at the end of Section 4 of [4]) is completely determined by the four mutually-cancelling comparisons $234 \succeq 15$, $1 \succeq 245$, $25 \succeq 34$, $45 \succeq 2$, and the extra conditions $235 \succ 1$ and $1 \succ 23$. Now if this ordering arises from the probability measure $\mathbf{p} = (p_1, \ldots, p_5)$, then the first four comparisons require that $p_1 = 2p_4 + 2p_5$, $p_2 = p_4 + p_5$ and $p_3 = 2p_5$, and the two extra conditions require that $2p_5 > p_4$ and $p_4 > p_5$. Choosing $p_4 = 3$ and $p_5 = 2$, for instance, we see that the given ordering almost agrees with the probability measure $\mathbf{p} = \frac{1}{24}(10, 5, 4, 3, 2)$.

Theorem 6. *There are exactly 169444 different comparative probability orderings of the subsets of $I_6 = \{1, 2, 3, 4, 5, 6\}$. Of these, 124187 are representable, while 45257 are not. Of the 45257 that are non-representable, 44987 fail C_4 with $a_1 = a_2 = a_3 = a_4 = 1$, while the other 270 satisfy C_4 but fail C_5 with $a_1 = a_2 = a_3 = a_4 = a_5 = 1$. Moreover, of these 45257 orderings that are non-representable, 40055 are almost representable, while the other 5202 are not. In particular, of the 44987 non-representable orderings that fail C_4, exactly 39785 are almost representable, while the other 5202 are not, and all of the 270 non-representable orderings that satisfy C_4 but fail C_5 are almost representable. Hence there are exactly 164242 almost representable orderings, and of these, 124187 are representable, while 40055 are not, with 39785 failing C_4 with $a_1 = a_2 = a_3 = a_4 = 1$, and the other 270 satisfying C_4 but failing C_5 with $a_1 = a_2 = a_3 = a_4 = a_5 = 1$.*

We say that the *dual* of a comparison $A \succ B$ is the comparison $B \succ A$, and then the dual of a set of comparisons is the set of their duals. In the case $n = 5$, Fishburn categorised the 30 orderings that are non-representable into five types (and their duals) according to the violations of C_4 that they exhibit, and noted that some orderings may exhibit more than one type of violation. For the case $n = 6$, we find that up to duality there are 423 such types, of which 385 (and their duals) involve four comparisons that together fail C_4 with $a_1 = a_2 = a_3 = a_4 = 1$, while the remaining 38 (and their duals) involve five comparisons that together fail C_5 with $a_1 = a_2 = a_3 = a_4 = a_5 = 1$.

We note that the first example of a comparative probability ordering on a 6-element set which fails to be almost representable was constructed by Kraft, Pratt, and Seidenberg [6]. Our computations provide the following extension to the results of Fishburn [4, 5]:

Corollary 3. $f(6) = g(6) = f^{\dagger}(6) = g^{\dagger}(6) = 5$.

6 Conclusion and Open Questions

Kraft, Pratt, and Seidenberg [6] left open the question of whether or not all comparative probability orderings on a 5-element set are almost representable. We have answered this question by showing that all of them are indeed almost representable.

We know now that $f(5) = g(5) = 4$, and $f(6) = g(6) = 5$, but $g(7) = 7$. The question remains open as to whether $f(7)$ is equal to 7 or 8. Nevertheless, it is clear that the situation for $n \geq 7$ is different to that for $n \in \{5, 6\}$, so Fishburn's conjecture is no longer adequate. We propose amending it to the following:

Conjecture: $f(n) = g(n) = n$ for all $n \geq 7$.

Some more research will have to be undertaken before any reasonable conjecture about $f^\dagger(n)$ and $g^\dagger(n)$ can be formulated.

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