Distinguishing Maps: Angles, Cliques, and Transpositions

Thomas W. Tucker

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Theorem (TWT *Elec J. Comb. 2011*). If $A \subset Aut(G)$ is no- τ and D(A, V) > 2, then G is $C_3, C_4, C_5, K_4, K_5, K_7, O_6, O_8$. For all except O_8 there is a unique (well-known) orientable map M with $Aut^+(M) = A$. There is no such map for O_8 .

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So $S_2(O_n) = O_{n+2}$ and $S_1(K_n) = K_{n+1}$.

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- 3. $K_{n,n}$ for n = 3, 4, 5
- 4. $O_6, O_8, C_3 \times C_3, K_{3,3,3}$ and the cube.

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Examples: Intransitive d < 3

Туре	map	G	surface	V-orb	E-orb			
bipart		$K_{1,n}, K_{2,n}$	sphere	2	1 or 2			
1 rad	Cn	$S_1(C_n)$	sphere	2	2			
2 rad	Cn	$S_2(C_n)$	sphere	3	3			
1 rad	T1	$S_1(K_4)$	torus	2	3			
1 rad	T2	$S_1(K_5)$	g = 2	2	3			
3 rad	K_4^P	$S_3(K_4)$	proj	2	2			
2 rad	T2	$S_2(K_5)$	g = 2	3	4			
bipart	B(m, n)	K _{m,n}	g = 3, 4, 6	2	1			
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3 rad	K_4^P	$S_3(K_4)$	proj	2	2			
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bipart	B(m, n)	K _{m,n}	g = 3, 4, 6	2	1			
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Here B(m, n) is map from $CM(Z_m \times Z_n, (a, a^{-1}b, b^{-1}))$, where we place a vertex in each a^m and b^n face, join through vertices of original graph (and throw out old vertices and edges). Can be viewed as hypermap.

Examples: Transitive, not $D^+(M) = 3$

Name	G	d	Stab	g	au
$CM(Z_4, (1, -1, 2))$	<i>K</i> 4	3	D_1	1	mix
$CM(Z_5, (1, -1, 2, -2))$	K_5	4	D_1	2	no
$CM(Z_2^3,(x,y,z))$	cube	3	<i>D</i> ₃	1	reg
$CM(Z_2^3, (x, x + y, y, y + z, z, z + x))$	<i>O</i> 8	6	<i>D</i> ₃	7	all
$CM(Q,(i,j,k)^b)$	<i>O</i> ₈	6	<i>D</i> ₃	5	no
$CM(Z_3^2,(x,y)^b)$	$C_3 \times C_3$	4	<i>D</i> ₄	1	reg
$CM(Z_3^2,(x,y,-x+y)^b)$	K _{3,3,3}	6	D_6	1	reg
$CM(Z_3^2,(x,x+y,y,y-x)^b)$	K_9	8	<i>D</i> ₄	10	all
B(3,3)	K _{3,3}	3	D _n	1	reg
B(4,4)	K _{4,4}	4	<i>D</i> ₄	3	reg
B(5,5)	K _{5,5}	5	D_5	6	reg

Also K_6 in projective plane.

Given $A \subset Aut(M)$ with an edge orbit E_A , we can "twist" all the edges in E_A , to get a new map (probably nonorientable) M' with Aut(M) = Aut(M') and action is the same on the vertex set (the underlying graph stays the same). Usual Petrie duality is to twist all edges.

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Example For standard map M for double pyramid $S_2(C_5)$ in the sphere, there are three edge-orbit using the D_5 action fixing the apexes. This map has D(M) = 3 under that action, so we can get 8 maps in all by applying various partial Petrie duals.

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So the above tables can lead to many different maps with same action of Aut(M) on V.

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Seiya and I have been not communicating well. He has article coming up in Disc. Math. attacking same problem with some partial results

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Finally, for c = 4 each K_4 has the symmetry of the tetrahedron and for c = 6 each K_6 has the symmetry of K_6 in the projective plane (antipodal quotient of the icosahedron).

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Remark Holds for graph G and $A \subset Aut(G)$ with natural dihedral vertex-stabilizers (i.e. A_v acts on d neighbors as D_d in the natural way).

An **angle** in a map is *uvw* where *uv* and *vw* are edges. If there is no edge *uw* then the angle is bf closed, otherwise **open**.

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Suppose that v has valence d. Then the measure of an angle, m(uvw), is one more than the number of intervening vertices in the cyclic order at v, either clockwise or counterclockwise, whichever is less. Thus $m(uvw) \le d/2$. Map automorphisms preserve angle measure.

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Since an angle reflector reserves orientation, this shows why in the orientation-preserving case there are no bent open angles, so you get K_n , O_{2n} , C_n .